A Convex Max-Flow Approach to Distribution Based Figure-Ground Separation

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Abstract. This study investigates a convex relaxation approach to figure-ground separation with a global distribution matching prior evaluated by the Bhattacharyya measure. The problem amounts to finding a region that most closely matches a known model distribution. It has been previously addressed by curve evolution, which leads to sub-optimal and computationally intensive algorithms, or by graph cuts, which result in metrication errors. Solving a sequence of convex sub-problems, the proposed relaxation is based on a novel bound of the Bhattacharyya measure which yields an algorithm robust to initial conditions. Furthermore, we propose a novel flow configuration that accounts for labeling-function variations, unlike existing configurations. This leads to a new max-flow formulation which is dual to the convex relaxed sub-problems we obtained. We further prove that such formulation yields exact and global solutions to the original, non-convex sub-problems. A comprehensive experimental evaluation on the Microsoft GrabCut database demonstrates that our approach yields improvements in optimality and accuracy over related recent methods.

Key words. image segmentation, convex optimization, max-flow min-cut, Bhattacharyya measure, distribution matching, iterative methods

1. Introduction. This study investigates the problem of partitioning an image into a figure and ground, so that the figure region is consistent with a known (a priori learned) distribution. The solution to this problem was recently shown very useful in the following computer vision tasks.

- Co-segmentation of image pairs: The problem consists of finding the same object (figure) in a pair of images and has attracted a significant research attention in recent years [3, 29, 36, 37, 47, 52, 51]. It has been shown very useful in object recognition/image retrieval [19, 26, 47, 48], image editing [3] and summarization [22];
Interactive image segmentation: using minimal user inputs, e.g., bounding boxes or scribbles, interactive segmentation embeds clues on user intention, thereby limiting the space of possible solutions and easing dramatically the problem. It has been intensively investigated in recent years [9, 34, 41, 46, 50].

Segmentation with offline learning: Segmenting groups of images that share similar properties occurs in important applications, e.g., in medical image analysis. In such scenarios, offline learning of prior information from segmented training images is very useful [6, 18].

Tracking: When tracking a target region within an image sequence, one can segment the current frame using information learned from the segmentation of a previous frame [5, 25, 56].

The problem amounts to optimizing a global measure of similarity/dissimilarity between distributions. Possible measures include the Bhattacharyya coefficient [4, 6, 5, 25, 41, 56], $L_p$ norm of the difference between histograms [37, 47, 51], Wasserstein distance [38], and Kullback-Leibler divergence [2]. However, the Bhattacharyya coefficient yielded outstanding performances recently [4, 41]. Furthermore, it has several desirable properties. First, it has a fixed (normalized) range which is convenient in many practical scenarios, e.g., in image retrieval and object recognition. Second, it is not affected by the choice of the bin width [1] (A proof of this property can be found in [1]). Third, it imposes a metric structure [20].

On the one hand, as several recent studies have shown [4, 6, 38, 35, 41, 47, 56, 25], global distribution-similarity measures can afford an accuracy unattainable with standard descriptions based on pixelwise information [9]. On the other hand, they generally lead to challenging optimization problems in both continuous or variational and discrete or graph-based settings.

In the continuous setting, the energy functional is defined over continuous variables and image domains [6, 25, 38, 56]. Active curve evolution driven by gradient-descent optimization has been the most prevalent and flexible choice in the literature to minimize global functionals [2, 6, 25, 56]. However, such choice suffers from local optima, slow convergence and high sensitivities to initialization. Complex implementation in numerics is also a significant limitation.
of active curve methods.

In the discrete setting, the problem is stated as a label assignment where the input images are considered as graphs defined over the positional arrays [9, 11, 12]. In this respect, graph cut techniques have been of intense interest recently in computer vision because they can guarantee global optima in nearly real time [9]. Unfortunately, global energies have been generally avoided because they are not directly amenable to graph cut optimization. The recent works in [4, 37, 41, 47] are notable exceptions which use relaxations via bounds or approximations of the energy so as to befit graph cut optimization. These works led to significant improvements over active curves in regard to optimality, speed, and accuracy [4]. However, the well-known grid bias (or metrification error) is a limitation of such graph-based approaches [42]. Reducing metric artifacts can be done by increasing the number of neighboring graph nodes, but this may result in a heavy computation and memory load [32]. In practice, it is well known that graph cut algorithms can yield a state-of-the-art empirical performance in the context of 2D grids with 4-neighborhood systems. They use heuristics that handle efficiently sparse grids [11]. However, the efficiency of the algorithms may decrease when moving from 2D to 3D grids or when using denser (larger neighborhood) grids.

Continuous convex relaxation approaches share the advantages of both active curves and graph cuts, and have recently attracted a significant research attention in image partitioning problems. Notable recent studies have shown the potential of convex optimization in solving the classical piecewise constant Mumford-Shah segmentation model [28, 39, 54] and its multiphase (multiregion) variant [15, 33, 55], as well as a histogram-based segmentation using the Wasserstein distance [38]. Moreover, unlike graph cuts, the implementation of convex relaxation approaches can be parallelized [55] to reduce computational time. The recent study by Klodt et al. [31] compared discrete and continuous optimization techniques in computer vision. It showed that, in the case of 3D or higher-dimensional grids, convex relaxation approaches outperform graph cuts in regard to speed and accuracy.
This study investigates an iterative convex relaxation approach to figure-ground separation with a global distribution-matching prior evaluated by the Bhattacharyya measure. Here following a summary of the main contributions:

- Solving a sequence of convex sub-problems, the proposed relaxation is based on a novel bound of the Bhattacharyya measure which can be viewed as a generalization of the bound we recently derived in our conference paper in [4]. The bound in [4] assumes that, at each iteration, the new figure region is enclosed within the region obtained at the previous iteration (refer to the illustration in Figure 3.1b), an assumption which requires the initial region to include the target region. The proposed bound relaxes such assumption (refer to the illustration in Figure 3.1a), yielding an algorithm which is more robust to initial conditions (refer to the examples in Figure 6.1).

- We propose a novel flow configuration (a simple 1-D example is depicted by the rightmost graph of Figure 2.1) which accounts for labeling-function variations, unlike existing configurations [53, 54] (refer the middle graph of Figure 2.1). This leads to a new max-flow formulation which is dual to the convex relaxed sub-problems we obtained. We further prove that such formulation yields exact and global solutions to the original, non-convex sub-problems.

- A comprehensive experimental evaluation on the Microsoft GrabCut database (50 images with ground truth labellings) demonstrates that our approach yields improvements in optimality and accuracy over the Bhattacharyya Measure Graph Cut (BMGC) method we recently presented in [4].

2. Preliminaries.

2.1. The Basic Max-Flow and Min-Cut Models. In this section, we review the dualities of max-flow and min-cut in both spatially discrete (graph) and continuous settings, based upon the classical graph and flow configurations where flows stream from the source node $s$ to each image node, and then from each image node to the sink (The leftmost and middle configurations in Figure 2.1 illustrate the discrete and continuous cases, respectively).
Figure 2.1. Flow Configurations. Left: spatially discrete (classical graph based) \[9\]; Middle: spatially continuous \[53, 54\]; Right: The proposed spatially continuous configuration.

Let $G = (\mathcal{V}, \mathcal{E})$ be a graph composed of a vertex set $\mathcal{V}$ and an edge set $\mathcal{E} \subset \mathcal{V} \times \mathcal{V}$. $\mathcal{V}$ includes all the nodes of the discretized image together with two additional terminals, a source $s$ and a sink $t$. $\mathcal{E}$ includes all the links of the neighborhood system of image nodes, e.g. 4 or 8 connected, as well as the links between the terminal $s$ or $t$ and each image node $u \in \mathcal{V}\{s, t\}$ (see the leftmost configuration in Figure 2.1).

For such a graph, an $s$-$t$ cut divides $\mathcal{V}$ into two disjoint sub-graphs:

$$\mathcal{V} = \mathcal{V}_s \cup \mathcal{V}_t, \quad \mathcal{V}_s \cap \mathcal{V}_t = \emptyset$$

with $\mathcal{V}_s$ (respectively $\mathcal{V}_t$) containing source $s$ (respectively sink $t$).

Let $C(e) \geq 0$ be the cost associated to each edge $e \in \mathcal{E}$. The energy of each $s$-$t$ cut is the sum of all edge costs $C(e)$ where $e \in \mathcal{E}_{st} \subset \mathcal{E}$ links two nodes, one in $\mathcal{V}_s$ and the other in $\mathcal{V}_t$.

On the one hand, the min-cut problem consists of finding a minimum-energy $s$-$t$ cut, i.e.,

$$\min_{\mathcal{E}_{st} \subset \mathcal{E}} \sum_{e \in \mathcal{E}_{st}} C(e). \quad (2.1)$$

On the other hand, one can regard each edge $e \in \mathcal{E}$ as a pipe and its edge cost $C(e)$ as the corresponding capacity. In this case, graph $G$ gives a flow ‘network’. Let $p(e)$ be the flow passing edge $e \in \mathcal{E}$. It is well-known that the min-cut problem in (2.1) amounts to finding the maximal flow from $s$ over the ‘network’ \[21, 24\], i.e.,

$$\max_{p_s} \sum_{v \in \mathcal{V}\{s, t\}} p_s(v) \quad (2.2)$$
subject to the corresponding flow capacity constraints, and flow conservation condition, i.e.,
the balance between in-coming and out-going flows at each node \( v \in V \setminus \{s, t\} \),
\[
\left[ \sum_{u \in \mathcal{N}(v)} p(e(u, v)) \right] - p_s(v) + p_t(v) = 0 \tag{2.3}
\]
where \( \mathcal{N}(v) \) denotes the neighborhood set of \( v \).

The recent studies by Yuan et al. [53, 54] proved that the equivalence between max-flow (2.2) and min-cut (2.1) holds in the spatially continuous context: given a continuous image domain \( \Omega \) and two terminals \( s \) and \( t \) linked to each \( x \in \Omega \), we have three types of flows (refer to the middle continuous configuration in Figure 2.1): the source flow \( p_s(x) \) from \( s \) to \( x \), the sink flow \( p_t(x) \) from \( x \) to \( t \), and the spatial flow \( p(x) \) for each \( x \in \Omega \). These three flow fields are constrained pointwise by properly defined capacities:
\[
p_s(x) \leq C_s(x), \quad p_t(x) \leq C_t(x), \quad |p(x)| \leq C(x) \tag{2.4}
\]
as well as a flow conservation condition:
\[
p_t(x) - p_s(x) + \text{div} \, p(x) = 0, \quad \forall x \in \Omega \tag{2.5}
\]
where divergence \( \text{div} \, p(x) \) evaluates the excess of the spatial flow around \( x \).

Finding the maximum total flow from \( s \) based on the above configuration yields the following continuous max-flow problem
\[
\max_{p_s, p_t, p} \int_\Omega p_s \, dx \quad \text{s.t.} \ (2.4) \text{ and } (2.5). \tag{2.6}
\]

Yuan et al. [53, 54] proved that the continuous max-flow formulation in (2.6) is equivalent to the continuous min-cut problem studied in [13, 39]:
\[
\min_{u(x) \in [0,1]} \int_\Omega (1 - u)C_s \, dx + \int_\Omega uC_t \, dx + \int_\Omega C \, |\nabla u| \, dx. \tag{2.7}
\]
Moreover, a continuous max-flow algorithm can be derived from (2.6) to compute efficiently and reliably a solution of (2.7) via convex optimization [8, 45].
It is worth noting that the pioneering work of Strang [49] is one of the first attempts to investigate the key idea of formulating min-cut problems in a continuous optimization framework. However, Strang’s max-flow formulation is different from the continuous min-cut formulation in (2.7) in the sense that it eventually results in a ratio min-cut.

2.2. The Bhattacharyya Distribution Matching Energy. Let $I : \Omega \subset \mathbb{R}^2 \to \mathcal{Z} \subset \mathbb{R}^n$ be an image function which maps domain $\Omega$ to a space $\mathcal{Z}$ of a photometric variable such as a color vector. The problem we tackle in this study consists of finding a region $R \subset \Omega$ whose distribution most closely matches a known reference distribution $\mathcal{M}$.

Let $\mathcal{P}$ be the nonparametric estimate of the distribution of $I$ within $R$:

$$
\mathcal{P}(z) = \frac{\int_R K_z(x) \, dx}{|R|}, \quad \forall z \in \mathcal{Z}
$$

(2.8)

with $|R|$ the area of region $R$, i.e., $|R| = \int_R dx$. Typical choices of $K(\cdot)$ include the Dirac function, which yields the histogram, or the Gaussian kernel [40]:

$$
K_z(x) = \frac{1}{(2\pi\sigma^2)^{(n/2)}} \exp \left( -\|z - I(x)\|^2 / 2\sigma^2 \right)
$$

(2.9)

where $\sigma$ the width of the kernel. Note that the Dirac function can be approximated by a Gaussian kernel with a very small value of $\sigma$.

Assume that the distribution of $I$ within the target object, denoted $\mathcal{M}$, is given. Let $B(\mathcal{P}, \mathcal{M})$ denotes the Bhattacharyya measure which evaluates the similarity between $\mathcal{P}$ and $\mathcal{M}$:

$$
B(\mathcal{P}, \mathcal{M}) = \sum_{z \in \mathcal{Z}} \sqrt{\mathcal{P}(z)\mathcal{M}(z)}
$$

(2.10)

Adding a regularization term for smooth region boundary, the problem amounts to the following optimization:

$$
\min_{\mathcal{R}} \left\{ E(\mathcal{R}) := -B(\mathcal{P}, \mathcal{M}) + C |\partial\mathcal{R}| \right\}
$$

(2.11)

where $|\partial\mathcal{R}|$ measures the boundary length of region $\mathcal{R}$, and $C$ a positive constant weighing the contribution of the regularization terms.
Let \( u : \Omega \to \{0, 1\} \) be the indicator function of \( R \subset \Omega \), i.e., \( u(x) = 1 \) when \( x \in R \) and \( u(x) = 0 \) otherwise. Using \( u \), the distribution in (2.8) can be rewritten as follows:

\[
\mathcal{P}(z) = \frac{\int_{\Omega} K_z(x) u(x) \, dx}{\int_{\Omega} u(x) \, dx}, \quad \forall z \in \mathcal{Z}
\]

(2.12)

and the Bhattacharyya similarity in (2.10) as follows:

\[
\mathcal{B}(u, \mathcal{M}) = \sum_{z \in \mathcal{Z}} \left( \frac{\int_{\Omega} T_{M,z}(x) u(x) \, dx}{\int_{\Omega} u(x) \, dx} \right)^{1/2}
\]

(2.13)

where \( T_{M,z}(x) = K_z(x) M(z) \).

The perimeter of region \( R \) can be evaluated by the total-variation of indicator function \( u \) [17, 27]:

\[
|\partial R| = \int_{\Omega} C |\nabla u(x)| \, dx
\]

(2.14)

In view of (2.13) and (2.14), we can, therefore, reformulate (2.11) in terms of the binary valued labeling function \( u \):

\[
\min_{u(x) \in \{0, 1\}} \left\{ E(u) := - \sum_{z \in \mathcal{Z}} \left( \frac{\int_{\Omega} T_{M,z}(x) u(x) \, dx}{\int_{\Omega} u(x) \, dx} \right)^{1/2} + \int_{\Omega} C |\nabla u(x)| \, dx \right\}.
\]

(2.15)

3. Global Minimization of Energy Upper Bound. Direct computation of (2.15) is a very challenging optimization problem, both in theory and numerics, due to the constraint of a pointwise binary valued variable and the high nonconvexity of the energy. (2.15) is nonconvex even when \( u \in \{0, 1\} \) is relaxed by \( u \in [0, 1] \), unlike the classical continuous min-cut problem in [39, 54].

We propose a novel iterative convex relaxation solution to (2.15), which globally minimizes
a sequence of upper bounds of $E(u)$, denoted $F(u, u^i)$, $i \geq 1$ ($i$ is the iteration number):

$$u^{i+1} = \min_{u \in \{0, 1\}} F(u, u^i), \quad i \geq 1 \quad \text{s.t.}$$

$$E(u) \leq F(u, u^i), \quad i \geq 1 \quad (3.1a)$$

$$E(u) = F(u, u) \quad \forall u : \Omega \to \{0, 1\} \quad (3.1c)$$

We will give the expression of $F(u, u^i)$ in the next section. Let us first show that, in general, optimizing any energy upper bound that follows the constraints in (3.1) yields a monotonically decreasing sequence of the original energy $E$ and, therefore, a minimum of $E$ at convergence.

**Proposition 3.1.** Sequence $E(u^i)$ corresponding to the sequence of minima obtained in (3.1) is monotonically decreasing, i.e., $E(u^{i+1}) \leq E(u^i)$.

**Proof.** Applying constraint (3.1c) to $u = u^i$ gives

$$E(u^i) = F(u^i, u^i) \quad (3.2)$$

Now, by definition of minimum in (3.1a), we have

$$F(u^{i+1}, u^i) \leq F(u', u^i) \quad (3.3)$$

combining (3.2) and (3.3) gives

$$F(u^{i+1}, u^i) \leq E(u^i) \quad (3.4)$$

Also, applying constraint (3.1b) to $u = u^{i+1}$ gives

$$E(u^{i+1}) \leq F(u^{i+1}, u^i) \quad (3.5)$$

Finally, combining (3.4) and (3.5) proves Prop. 3.1.

**Proposition 3.2.** Sequence $E(u^i)$ is convergent.

**Proof.** $E(u^i)$ is monotonically decreasing and lower bounded, because the Bhattacharyya measure is upper bounded by one. Therefore $E(u^i)$ converges.
3.1. Energy Upper Bounds. In this section, we derive an energy upper bound $F(u, u^i)$, which is convex with respect to labeling function $u$. We will further show that $F(u, u^i)$ can be minimized globally and exactly over $u \in \{0, 1\}$ in the next sections.

Let $u^i \in \{0, 1\}$ denote the binary labeling function, obtained at the previous $i$-th step, as the starting point of the $i + 1$-th iteration. At this stage, we are seeking an optimal labeling function $u \in \{0, 1\}$ which minimizes the energy upper bound. Let us express the difference between $u$ and the given $u^i \in \{0, 1\}$ as a function of two new variables $u^+ \in \{0, 1\}$ and $u^- \in \{0, 1\}$ (refer to Figure 3.1a):

1. $u^+$ indicates the area where $u^i$ is 0 and $u$ becomes 1, i.e., area increase:

   $$u^+(x) := \begin{cases} u(x), & \text{where } u^i(x) = 0 \\ 0, & \text{otherwise} \end{cases}$$

   (3.6)

2. $u^-$ indicates the area where $u^i$ is 1 and $u$ becomes 0, i.e., area decrease:

   $$u^-(x) := \begin{cases} 1 - u(x), & \text{where } u^i(x) = 1 \\ 0, & \text{otherwise} \end{cases}$$

   (3.7)

The Bhattacharyya measure (2.13) can then be rewritten in terms of $u^+$ and $u^-$ as follows:

$$B(u, \mathcal{M}) = \sum_{z \in \mathcal{Z}} \left( \frac{\int_{\Omega}(\mathcal{T}_{M,z}u^i + \mathcal{T}_{M,z}u^+ - \mathcal{T}_{M,z}u^-) \, dx}{\int_{\Omega}(u^i + u^+ - u^-) \, dx} \right)^{1/2}.$$  

(3.8)
The following describes the upper bound we propose.

**Proposition 3.3.** Given a labeling $u^i \in \{0, 1\}$, for any labeling function $u \in \{0, 1\}$, we have the following upper bound which depends on $u^+$ and $u^-$:

$$E(u) \leq -B(u^i, \mathcal{M}) + J(u, u^i)$$  \hspace{1cm} (3.9)

where

$$J(u, u^i) = \int_{\Omega} C_v^i u^- \, dx + \int_{\Omega} C_w^i u^+ \, dx + \int_{\Omega} C |\nabla u| \, dx$$  \hspace{1cm} (3.10)

and

$$C_v^i(x) = \sum_{z \in \mathcal{Z}} \frac{D_{z,i} T_{M,z}(x)}{\int_{\Omega} T_{M,z}(x) u^i(x) \, dx},$$  \hspace{1cm} (3.11)

$$C_w^i(x) = \sum_{z \in \mathcal{Z}} \frac{D_{z,i}^2}{2 \int_{\Omega} u^i(x) \, dx},$$  \hspace{1cm} (3.12)

$$D_{z,i} = \left( \frac{\int_{\Omega} T_{M,z}(x) u^i(x) \, dx}{\int_{\Omega} u^i(x) \, dx} \right)^{1/2}.$$  \hspace{1cm} (3.13)

**Proof.** We define

$$v^+ = \frac{\int_{\Omega} T_{M,z}(x) u^+(x) \, dx}{\int_{\Omega} T_{M,z}(x) u^i(x) \, dx}, \quad v^- = \frac{\int_{\Omega} T_{M,z}(x) u^-(x) \, dx}{\int_{\Omega} T_{M,z}(x) u^i(x) \, dx};$$

and

$$w^+ = \frac{\int_{\Omega} u^+(x) \, dx}{\int_{\Omega} u^i(x) \, dx}, \quad w^- = \frac{\int_{\Omega} u^-(x) \, dx}{\int_{\Omega} u^i(x) \, dx}.$$  

Therefore, the Bhattacharyya measure (3.8) can be equally written as

$$B(\mathcal{P}, \mathcal{M}) = \sum_{z \in \mathcal{Z}} \left\{ \left( \frac{\int_{\Omega} T_{M,z}(x) u^i(x) \, dx}{\int_{\Omega} u^i(x) \, dx} \right)^{1/2} \frac{1 + v^+ - v^-}{\sqrt{1 + w^+ - w^-}} \right\}. \hspace{1cm} (3.14)$$
Observe all the variables $v^+$, $v^-$, $w^+$ and $w^-$ are positive, we have

$$ B(u, \mathcal{M}) \geq \sum_{z \in \mathcal{Z}} \left\{ \left( \frac{\int_{\Omega} T_{M,z}(x) u^i(x) \, dx}{\int_{\Omega} u^i(x) \, dx} \right)^{1/2} \sqrt{\frac{1-v^-}{1+w^+}} \right\} \quad (3.15) $$

$$ \geq \sum_{z \in \mathcal{Z}} \left\{ \left( \frac{\int_{\Omega} T_{M,z}(x) u^i(x) \, dx}{\int_{\Omega} u^i(x) \, dx} \right)^{1/2} \left( 1-v^- \right) \left( 1 - \frac{1}{2} w^+ \right) \right\} \quad (3.16) $$

$$ \geq \sum_{z \in \mathcal{Z}} \left\{ \left( \frac{\int_{\Omega} T_{M,z}(x) u^i(x) \, dx}{\int_{\Omega} u^i(x) \, dx} \right)^{1/2} \left( 1-v^- - \frac{1}{2} w^+ \right) \right\}. \quad (3.17) $$

The above result (3.16) can be seen by the facts that for any $v^- \in [0,1]$ and $w^+ \in \mathbb{R}^+$, we have

$$ \sqrt{1-v^-} \geq 1 - v^-, \quad 1/\sqrt{1+w^+} \geq 1 - \frac{1}{2} w^+. $$

In view of (3.17), we have

$$ B(u, \mathcal{M}) \geq B(u^i, \mathcal{M}) - \int_{\Omega} C^i_v(x) u^-(x) \, dx - \int_{\Omega} C^i_w(x) u^+(x) \, dx \quad (3.18) $$

where

$$ C^i_v(x) = \sum_{z \in \mathcal{Z}} \frac{\mathcal{D}_{z,n} T_{M,z}(x)}{\int_{\Omega} T_{M,z}(x) u^i(x) \, dx}, \quad C^i_w(x) = \sum_{z \in \mathcal{Z}} \frac{\mathcal{D}_{z,n}}{2 \int_{\Omega} u^i(x) \, dx} $$

and

$$ \mathcal{D}_{z,n} = \left( \frac{\int_{\Omega} T_{M,z}(x) u^i(x) \, dx}{\int_{\Omega} u^i(x) \, dx} \right)^{1/2}. $$

Therefore, by (3.18), Prop. 3.3 can be proved.

Prop. 3.3 leads us to the following conclusions:

- **Energy upper bound:** In view of (3.9), we have the energy upper bound

$$ F(u, u^i) = -B(u^i, \mathcal{M}) + J(u, u^i) \quad (3.19) $$

where $B(u^i, \mathcal{M})$ is a constant corresponding to the given $u^i$.

- **Continuous min-cut model:** Hence we can minimize $F(u, u^i)$ over $u \in \{0,1\}$ by

$$ \min_{u \in \{0,1\}} J(u, u^i) \quad (3.20) $$
where $J(u, u^i)$ is given in (3.10). The bound in (3.19) and its variable part $J(u, u^i)$ do not reference labeling variable $u$ in the same way the standard min-cut problem in (2.7) does. $J(u, u^i)$ depends rather on labeling-change variables ($u^+$ and $u^-$). Therefore, it is not straightforward to optimize $J(u, u^i)$ with the standard flow configurations in the left and middle parts of Figure 2.1. The proposed flow configuration in the right part of Figure 2.1 is designed so as to account for $u^+$ and $u^-$. In the sequel of this paper, we will call (3.20) the continuous min-cut model.

- **Cost of changing labels:** It is easy to see that $C_{iv}^i(x)$ evaluates the cost of changing label $u(x)$ from $u^i(x) = 1$ to $u(x) = 0$, i.e., $x$ is previously labeled as foreground and cost $C_{iv}^i(x)$ is paid when $x$ is relabeled as background. $C_{iw}^i(x)$ evaluates the cost of moving $x$ from the background to the foreground.

4. **Continuous Max-Flow Approach.** The new continuous min-cut model in (3.20) depends on labeling-function variations ($u^+$ and $u^-$), not on a single labeling function as with previous models [13, 39, 53, 54]. First, we examine a convex relaxation of (3.20). Then, we propose a novel flow configuration (refer to the rightmost graph of Figure 2.1) to solve such relaxation. This leads to a new max-flow formulation which is dual to the convex relaxed version of (3.20). Finally, we prove that our formulation yields a global and exact solution of (3.20).

4.1. **Convex Relaxed Min-Cut Model.** Let $\Omega_s$ and $\Omega_t$ two disjoint domains given by the current labeling function $u^i \in \{0, 1\}$:

$$\Omega = \Omega_s \cup \Omega_t, \quad \Omega_s \cap \Omega_t = \emptyset$$

(4.1)

where $\Omega_s$ is indicated by $u^i = 1$ and $\Omega_t$ by $u^i = 0$.

From the definitions of $u^+$ (3.6) and $u^-$ (3.7), we reformulate (3.20) by:

$$\min_{u(x) \in \{0, 1\}} \int_{\Omega_t} C_{iv}^i (1-u) \, dx + \int_{\Omega_s} C_{iw}^i u \, dx + \int_{\Omega} C |\nabla u| \, dx$$

(4.2)
and propose a convex relaxation of (4.2) as follows:

$$\min_{u(x) \in [0,1]} \int_{\Omega_s} C^v u (1 - u) \, dx + \int_{\Omega_t} C^w u \, dx + \int_{\Omega} C |\nabla u| \, dx$$  \hspace{1cm} (4.3)$$

where the binary-valued constraint $u \in \{0,1\}$ is relaxed to $u \in [0,1]$. In the remainder of the paper, we will call (4.3) the convex relaxed min-cut model.

### 4.2. Continuous Max-Flow Model

We propose a novel configuration of flows (a simple 1-D example is depicted by the rightmost graph of Figure 2.1) which accounts for the previous partition defined by $\Omega_s$ and $\Omega_t$, unlike the existing continuous flow setting [53, 54] (refer the middle graph of Figure 2.1). Given $\Omega_s$ and $\Omega_t$, we define three types of flows: the source flow $p_s(x)$ directed from $s$ to $\forall x \in \Omega_s$, the sink flow $p_t(x)$ directed from $\forall x \in \Omega_t$ to $t$, and the spatial flow $p(x)$ given $\forall x \in \Omega$.

$p_s(x)$, $p_t(x)$ and $p(x)$ are constrained by the following flow capacities:

$$p_s(x) \leq C^v u(x), \quad \forall x \in \Omega_s; \quad \text{(4.4)}$$

$$p_t(x) \leq C^w u(x), \quad \forall x \in \Omega_t; \quad \text{(4.5)}$$

$$|p(x)| \leq C, \quad \forall x \in \Omega; \quad \text{(4.6)}$$

and flow conservation conditions:

$$-p_s(x) + \text{div} \, p(x) = 0, \quad \forall x \in \Omega_s; \quad \text{(4.7)}$$

$$p_t(x) + \text{div} \, p(x) = 0, \quad \forall x \in \Omega_t. \quad \text{(4.8)}$$

The continuous max-flow model can, therefore, be formulated by maximizing the total flow from the source $s$:

$$\max_{p_s,p_t,p} \int_{\Omega_s} p_s \, dx \quad \text{s.t. (4.4) - (4.8)} \quad \text{(4.9)}$$

It is worth noting that the maximum in (4.9) is attained because it is a linear problem subject to linear constraints.
4.3. Equivalence between Continuous Max-Flow (4.9) and Continuous Min-Cut (4.3).

Introduce the multiplier function $u(x) \in \mathbb{R}$ to the flow conservation constraints (4.7) at $\forall x \in \Omega_s$ and (4.8) at $\forall x \in \Omega_t$. We have the following primal-dual formulation equivalent to (4.9):

$$\min_u \max_{p_s,p_t,p} \int_{\Omega_s} p_s \, dx + \int_{\Omega_s} (\text{div} \, p - p_s)u \, dx + \int_{\Omega_t} (\text{div} \, p + p_t)u \, dx \quad \text{s.t.} \ (4.4) - (4.6)$$

which can be rearranged as

$$\min_u \max_{p_s,p_t,p} \int_{\Omega_s} (1-u)p_s \, dx + \int_{\Omega_t} up_t \, dx + \int_\Omega u \, \text{div} \, p \, dx \quad \text{s.t.} \ (4.4) - (4.6). \quad (4.11)$$

To build the equivalence between max-flow model (4.9) and min-cut formulation (4.3), we prove

**Proposition 4.1.** The primal-dual model (4.11) amounts to

$$\min_u \int_{\Omega_s} (1-u)C_v \, dx + \int_{\Omega_t} uC_w \, dx + \int_\Omega C \, |\nabla u| \, dx \quad (4.12)$$

where

$$u(x) \leq 1, \quad \forall x \in \Omega_s; \quad u(x) \geq 0, \quad \forall x \in \Omega_t.$$

**Proof.** The proof is based on the following variational facts:

The constraints of flows are convex, and the energy function is linear in both the primal and dual functions $p_s(x), p_t(x), p(x)$ and $u(x)$, hence convex l.s.c. (lower semi-continuous) for fixed $u$ and concave u.s.c. (upper semi-continuous) for fixed $p_s(x), p_t(x)$ and $p(x)$ [45]. This implies the existence of at least one saddle point of the primal-dual model (4.10) or (4.11), i.e., strong duality. It follows that we can interchange the min and max order of (4.10) or (4.11). Here we first maximize the energy function of the primal-dual model over the three flow functions $p_s(x), p_t(x), p(x)$ and then minimize over $u(x)$.

Now we consider the optimization problem

$$f(v) = \sup_{w \leq C} vw, \quad (4.13)$$
where \(v\), \(w\) and \(C\) are scalars. When \(v < 0\), \(w\) can be negative infinity in order to maximize the value \(vw\), i.e., \(f(v) = +\infty\). To achieve a meaningful function \(f(v)\) over \(v\), we have \(v \leq 0\).

It can also be easily seen that for \(w \leq C\),

\[
\begin{cases}
  \text{if } v = 0, & \text{then } w \leq C \text{ maximizes the function and } f(v) = 0, \\
  \text{if } v > 0, & \text{then } w = C \text{ maximizes the function and } f(v) = vC.
\end{cases}
\]  

(4.14)

In summary, we have

\[f(v) = vC, \quad \text{and } v \geq 0.\]  

(4.15)

In view of (4.13) and (4.15), it follows that the maximization of (4.11) over \(p_s\), constrained by (4.4), amounts to

\[
\max_{p_s(x) \leq C^i_s(x)} \int_{\Omega_s} (1 - u(x)) p_s(x) \, dx = \int_{\Omega_s} (1 - u(x)) C^i_s u(x) \, dx, \quad \text{and } u(x) \leq 1, \text{ for } x \in \Omega_s.
\]  

(4.16)

In the same manner, the maximization of (4.11) over \(p_t\), constrained by (4.5), corresponds to

\[
\max_{p_t(x) \leq C^i_t(x)} \int_{\Omega_t} u(x) p_t(x) \, dx = \int_{\Omega_t} u(x) C^i_t w(x) \, dx, \quad \text{and } u(x) \geq 0, \text{ for } x \in \Omega_t.
\]  

(4.17)

It is also well-known that [27]:

\[
\max_{|p(x)| \leq C(x)} \int_{\Omega} u \div p \, dx = \int_{\Omega} C |\nabla u| \, dx.
\]  

(4.18)

Observe (4.16), (4.17) and (4.18), Prop. 4.1 is proved. \(\blacksquare\)

Observe Prop. 4.1, we further prove

Proposition 4.2. The optimum \(u^*\) of the minimization problem in (4.12) verifies \(u^* \in [0, 1]\) for almost everywhere in \(\Omega\). Therefore, we have the equivalence:

\[(4.9) \iff (4.10) \text{ or } (4.11) \iff (4.3).\]

Proof. Let \(u^*\) be the minimum of (4.12). Due to the convexity of (4.12), \(u^*\) is simply accepted as a global minimum. Now we prove that \(0 \leq u^*(x) \leq 1 \forall x \in \Omega\).
Through Prop. 4.1, we have

\[ u^*(x) \leq 1, \quad \text{for } x \in \Omega_s; \quad u^*(x) \geq 0, \quad \text{for } x \in \Omega_t. \]

Now we prove \( u^*(x) \geq 0 \), for \( x \in \Omega_s \). If \( u^*(x) < 0 \) at \( \tilde{\Omega} \subset \Omega_s \), then we define the function \( u'(x) \) which just thresholds the value \( u^*(x) \) to be non-negative, i.e.,

\[
u'(x) = \begin{cases} 
0 & \text{for } x \in \tilde{\Omega} \\
u^*(x) & \text{for } x \in \Omega \setminus \tilde{\Omega}.
\end{cases}
\]

Obviously \( u'(x) \leq u^*(x) \), for \( x \in \Omega_s \), we have

\[
\int_{\Omega_s} (1 - u^*(x)) C^i_v(x) \, dx = \int_{\Omega_s} (1 - u'(x)) C^i_v(x) \, dx + \int_{\tilde{\Omega}} (0 - u^*(x)) C^i_v(x) \, dx.
\]

Observe that \( C^i_v(x) > 0 \) by (3.11) and \( u^*(x) < 0 \) at \( \tilde{\Omega} \subset \Omega_s \), we have

\[
\int_{\Omega_s} (1 - u'(x)) C^i_v(x) \, dx < \int_{\Omega_s} (1 - u^*(x)) C^i_v(x) \, dx \quad (4.19)
\]

On the other hand, by the coarea formula of the total variation term:

\[
\int_{\Omega} C |\nabla u| \, dx = \int_{-\infty}^{+\infty} L_\gamma(u) \, d\gamma,
\]

where \( L_\gamma(u) \) is the weighted length of the \( \gamma \)-upper level set of \( u(x) \). It follows that

\[
\int_{\Omega} C |\nabla u'| \, dx < \int_{\Omega} C |\nabla u^*| \, dx, \quad (4.20)
\]

because for \( u'(x) \), all the \( \gamma \)-upper level sets of \( u' \), when \( \gamma < 0 \), are thresholded to vanish.

Through (4.19), (4.20) and \( u'(x) = u^*(x) \) for \( x \in \Omega_t \), we have

\[
\left\{ \int_{\Omega} (1 - u'(x)) C^i_v(x) \, dx + \int_{\Omega_t} u'(x) C^i_w(x) \, dx + \int_{\tilde{\Omega}} C |\nabla u'| \, dx \right\} < \left\{ \int_{\Omega_s} (1 - u^*(x)) C^i_v(x) \, dx + \int_{\Omega_t} u^*(x) C^i_w(x) \, dx + \int_{\tilde{\Omega}} C |\nabla u^*| \, dx \right\}.
\]

This is in contradiction to the fact that \( u^*(x) \) is the global optimum of (4.12).

Likewise, we can prove that \( u^*(x) \leq 1 \), for \( x \in \Omega_t \). We, therefore, have \( u^*(x) \in [0, 1] \) for \( x \in \Omega \).

The equivalence between the continuous max-flow model (4.9), the primal-dual model (4.10) or (4.11) and the dual convex min-cut model (4.3) can be established consequently.
4.4. Exactness of the Convex Relaxed Min-Cut Model in (4.3). In this section, we show that the binary-valued optimization problem in (4.2) can be solved globally and exactly by its convex relaxation in (4.3).

Clearly, (4.3) is convex. Let \( u^*(x) \) be its global optimum. For any \( \gamma \in [0, 1] \), we define the \( \gamma \)-upper level set \( u^*_\gamma(x) \) for all \( x \in \Omega \) by

\[
u^*_\gamma(x) = \begin{cases} 
1, & \text{when } u^*(x) > \gamma \\
0, & \text{when } u^*(x) \leq \gamma
\end{cases} \tag{4.21}
\]

Moreover, we have

**Proposition 4.3.** When \( u^*(x) \) gives one global optimum of (4.3), its thresholding \( u^*_\gamma(x) \in \{0, 1\} \) by (4.21), for any \( \gamma \in [0, 1] \), solves the binary-valued optimization problem in (4.2) globally and exactly.

The proof of Prop. 4.3 follows the ideas of [53, 54], is discussed in [14, 16, 39] and is given in below.

**Proof.** There exists \((p^*_s, p^*_t, p^*; u^*(x))\) to be the optimal primal-dual pair of the primal-dual model (4.10) or (4.11), see the proof of Prop. 4.1. Clearly, \( p^*_s, p^*_t, p^* \) optimize the continuous max-flow problem (4.9) and gives the globally maximal max-flow energy

\[ E_{mf}^* = \int_{\Omega} p^*_s(x) \, dx. \]

\( u^*(x) \) optimizes the dual convex min-cut formulation (4.3) and gives the globally minimal energy of the convex min-cut model (4.3), i.e.,

\[ E_{mc}^* = \int_{\Omega_s} (1 - u^*(x)) C^i_v(x) \, dx + \int_{\Omega_t} u^*(x) C^i_w(x) \, dx + \int_{\Omega} C |\nabla u^*| \, dx. \]

Let \( E_{pd}(p_s, p_t, p; u) \) the primal-dual energy of (4.10) over the minimax of \( p_s, p_t, p \) and \( u \), i.e.,

\[ E_{pd}(p_s, p_t, p; u) = \int_{\Omega_s} (1 - u(x)) p_s(x) \, dx + \int_{\Omega_t} u p_t(x) \, dx + \int_{\Omega} u(x) \text{ div } p(x) \, dx. \]
Observe the proof of Prop. 4.1 (the interchangeable min-max order), we have

$$E_{mf}^* = \max_{p_s, p_t, p} \min_u E_{pd}(p_s, p_t, p; u) = E_{pd}(p_s^*, p_t^*, p^*; u^*) = \min_{u} \max_{p_s, p_t, p} E_{pd}(p_s, p_t, p; u) = E_{mc}^*. \quad (4.22)$$

Given the $\gamma$-level set $u_\gamma^*(x)$ of $u^*(x)$, for any $\gamma \in (0, 1]$, now we prove $(p_s^*, p_t^*, p^*; u_\gamma^*)$ is also an optimal primal-dual pair. This can be validated by the following facts:

- When the flow function $p_s^*(x)$ maximizes $\int_{\Omega} u \text{div} p_s^* \, dx$, then for any $\gamma$-upper level set of $u(x)$, i.e., $u_\gamma(x)$, $\gamma \in (u_{\text{min}}, u_{\text{max}}]$, $p^*(x)$ also maximizes $\int_{\Omega} u_\gamma \text{div} p^* \, dx$. This is another representation of the coarea theorem. Then, it follows that $p^*(x)$ also maximizes integral $\int_{\Omega} u_\gamma^* \text{div} p^* \, dx$.

- Let $\tilde{\Omega} \subset \Omega_s$ be the area where $u^*(x) < \gamma$ for $x \in \tilde{\Omega}$. In $\Omega_s$, $p_s^*$ maximizes the integral $\int_{\Omega_s} (1 - u^*) \, p_s \, dx$ over the flow capacity constraint (4.4) for $p_s$. Observe (4.14) and $1 - u^*(x) > 0$ for $x \in \tilde{\Omega}$, we have $p_s^*(x) = C_s(x)$ for $x \in \tilde{\Omega}$. In addition, for $u_\gamma^*(x)$, we have
  
  $$1 - u_\gamma^*(x) = 1, \quad \text{for } x \in \tilde{\Omega}; \quad 1 - u_\gamma^*(x) = 0, \quad \text{for } x \in \Omega_s \setminus \tilde{\Omega}. $$

  Therefore, by (4.14), $p_s^*(x)$ maximizes the integral $\int_{\Omega_s} (1 - u^*_\gamma) \, p_s \, dx$ over the flow capacity constraint (4.4).

- Let $\tilde{\Omega} \subset \Omega_t$ be the area where $u^*(x) \geq \gamma$ for $x \in \tilde{\Omega}$. In $\Omega_t$, $p_t^*$ maximizes the integral $\int_{\Omega_t} u^* \, p_t \, dx$ over the flow capacity constraint (4.5) for $p_t$. Observe (4.14) and $u^*(x) > 0$ for $x \in \tilde{\Omega}$, we have $p_t^*(x) = C_t(x)$ for $x \in \tilde{\Omega}$. In addition, for $u_\gamma^*(x)$, we have
  
  $$u_\gamma^*(x) = 1, \quad \text{for } x \in \tilde{\Omega}; \quad u_\gamma^*(x) = 0, \quad \text{for } x \in \Omega_t \setminus \tilde{\Omega}. $$

  Therefore, by (4.14), $p_t^*(x)$ maximizes the integral $\int_{\Omega_t} u^*_\gamma \, p_t \, dx$ over the flow capacity constraint (4.4).
Hence, we have

\[
\max_{p_s, p_t, p} E_{pd}(p_s, p_t, p; u^*_\gamma) = E_{pd}(p^*_s, p^*_t, p^*; u^*_\gamma)
\]

\[
= \int_{\Omega_s} p^*_s(x) \, dx + \int_{\Omega_t} (\text{div } p^* - p^*_s) u^*_\gamma \, dx + \int_{\Omega_t} (\text{div } p^* + p^*_t) u^*_\gamma \, dx
\]

\[
= \int_{\Omega_s} p^*_s(x) \, dx = E^{mf}_{\text{opt}}
\]

(4.23)

where the last term comes from the flow conservation conditions (4.7) and (4.8).

On the other hand, we also have

\[
E_{pd}(p^*_s, p^*_t, p^*; u^*_\gamma) = \int_{\Omega_s} (1 - u^*_\gamma(x)) \, C_v^i(x) \, dx + \int_{\Omega_t} u^*_\gamma(x) \, C_w^i(x) \, dx + \int_{\Omega} C |\nabla u^*_\gamma| \, dx ,
\]

(4.24)

following the same procedure of the proof for Prop.4.

Combining (4.22), (4.23) and (4.24), we have

\[
\int_{\Omega_s} (1 - u^*_\gamma(x)) \, C_v^i(x) \, dx + \int_{\Omega_t} u^*_\gamma(x) \, C_w^i(x) \, dx + \int_{\Omega} C |\nabla u^*_\gamma| \, dx
\]

\[
= \int_{\Omega_s} (1 - u^*(x)) \, C_v^i(x) \, dx + \int_{\Omega_t} u^*(x) \, C_w^i(x) \, dx + \int_{\Omega} C |\nabla u^*_\gamma| \, dx
\]

\[
= E^{me}_{\text{opt}}.
\]

Observing \(E^{me}_{\text{opt}}\) gives the global minimal energy of the convex min-cut (4.3), the relaxation of (4.3), we conclude that \(u^*_\gamma(x) \in \{0, 1\}\) solves globally and exactly (4.2).

5. Continuous Max-Flow Based Algorithm. The following summarizes the proposed procedure:

1. Start with an arbitrary initial labeling \(u^i (i = 1)\);
2. For the \(i\)-th outer iteration, set the flow configuration based on labeling function \(u^i\).
   Compute \(C_v^i\) and \(C_w^i\) by (3.11) and (3.12);
3. Compute \(u^{i+1}(x) \in \{0, 1\}\) as the global and exact optimum in (4.2) via the convex relaxation in (4.3) and the corresponding max-flow model in (4.9). The max-flow algorithm is detailed in Alg. 1;
4. Let \(i = i + 1\) and repeat the above two steps until convergence.
Note that the energy in (4.10) is the Lagrangian function of (4.9), where labeling $u$ is the multiplier to flow conservation conditions (4.7) and (4.8). In this regard, we define the following augmented Lagrangian function:

$$
L_c(p_s, p_t, p, u) = \int_{\Omega_s} p_s \, dx + \int_{\Omega_s} (\text{div} \, p - p_s) u \, dx \\
+ \int_{\Omega_t} (\text{div} \, p + p_t) u \, dx - \frac{c}{2} \|\text{div} \, p - p_s\|_{\Omega_s}^2 \\
- \frac{c}{2} \|\text{div} \, p + p_t\|_{\Omega_t}^2
$$

where $c > 0$. Finally, we derive the continuous max-flow algorithm from the classical augmented Lagrangian method [43, 44], whose convergence properties can be found in [7, 8]. The solution is sought by optimization over the flows and labeling simultaneously.

**Algorithm 1 Continuous Max-Flow Based Algorithm**

- $i = 1$: Initialize flows $p_s^1, p_t^1, p^1$ and labeling $u^1 \in \{0, 1\}$. Then start $i + 1$-th iteration;

- Fix $p^i$ and $u$ and maximize in a closed form $L_c(p_s, p_t, p, u)$ over $p_s$ and $p_t$:

$$
(p_s^{i+1}, p_t^{i+1}) := \arg \max_{p_s, p_t} L_c(p_s, p_t, p^i, u);
$$

- Fix $p_s^{i+1}, p_t^{i+1}$ and $u$ and maximize iteratively $L_c(p_s, p_t, p, u)$ over $p$ via the simple projection-descent step:

$$
p^{i+1} := \arg \max_p L_c(p_s^{i+1}, p_t^{i+1}, p, u);
$$

- Update labeling function $u$ by

$$
u(x) = \begin{cases} 
    p(x) + c (\text{div} \, p^{i+1}(x) - p_s^{i+1}(x)), & \forall x \in \Omega_s \\
    p(x) + c (\text{div} \, p^{i+1}(x) + p_t^{i+1}(x)), & \forall x \in \Omega_t
\end{cases}
$$

Repeat the above two steps until convergence.

6. **Experiments.** We applied the proposed method to the 50 images of the Microsoft GrabCut segmentation database [47], which also contains ground truth segmentations of the foreground objects. The results were compared with similar experiments on the same data.
using several other algorithms. We measured the optimality by computing (1) the Bhattacharyya measure at convergence and (2) the difference in histogram w.r.t. ground truth, i.e., the $L_1$ norm of the difference between the obtained region and ground truth histograms. We also evaluated the accuracy of the algorithms by computing the number of misclassified pixels in comparison to the ground truth. For the proposed algorithm, the parameters for all 50 images were selected as follows: $C = 0.3$, error threshold (convergence criteria) = $1 \times 10^{-2}$ and $c = 0.05$. The density is estimated from the RGB components with a 3-dimensional histogram consisting of $256 \times 256 \times 256$ bins. Following the experiment in [47], we estimate the model from the ground truth for all compared methods.

6.1. Quantitative evaluations: Table 6.1 reports average number of misclassified pixels (error), the norm of the average difference between the obtained histogram and the model ($L_1$-norm), average computational time (time), the Bhattacharyya measure at convergence ($B$) and the number of KDE (Kernel Density Estimation) evaluations per image (# KDEs) for:

- The proposed method;
- The bound in [4] in conjunction with the standard convex max-flow framework in [53, 54];
- The bound in [4] optimized via discrete Graph Cuts;
- The log-likelihood model in [10, 46] optimized via discrete Graph Cuts; and
- Active curve optimization [6, 25];

The proposed method yielded the lowest error, highest Bhattacharyya measure and lowest $L_1$-norm of histogram difference, which indicates improvements in accuracy and optimality. For the proposed method, the average time for computing the kernel density estimates and the Bhattacharyya distances is 4.5 seconds, which corresponds to approximately 25% of the overall computational time (18.8 seconds).

Note that the figure and ground prior distributions required by the likelihood-based model [10, 46] were learned from the ground truth and, therefore, known in advance. Even though
### Methods

<table>
<thead>
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<th>Methods</th>
<th>$L_1$-norm</th>
<th>$\mathcal{B}$</th>
<th># KDEs</th>
<th>time</th>
<th>error</th>
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<tbody>
<tr>
<td>Proposed flow config. + proposed bound</td>
<td>373.2</td>
<td>0.9991 ± 0.0009</td>
<td>13</td>
<td>18.8</td>
<td>383</td>
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<tr>
<td>Flow config. in [53, 54] + bound in [4]</td>
<td>737.5</td>
<td>0.9943 ± 0.0059</td>
<td>3</td>
<td>3.2</td>
<td>750</td>
</tr>
<tr>
<td>Graph Cuts + bound in [4]</td>
<td>501.1</td>
<td>0.9984 ± 0.0016</td>
<td>4</td>
<td>3.5</td>
<td>520</td>
</tr>
<tr>
<td>Log-likelihood model in [10, 46]</td>
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<td>0.9903 ± 0.0071</td>
<td>1</td>
<td>0.2</td>
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<tr>
<td>Active curve optimization</td>
<td>-</td>
<td>0.9791 ± 0.01</td>
<td>1737</td>
<td>631</td>
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<table>
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<th>Table 6.1</th>
</tr>
</thead>
<tbody>
<tr>
<td>Quantitative accuracy and optimality evaluations:</td>
</tr>
</tbody>
</table>

this comparison is not in favor of our method (as it does not use the background-model information), our experiments indicate that the proposed algorithm can bring an improvement of the order of $10^3$ pixels.

**6.2. Robustness w.r.t. initial conditions.** The examples in Figure 6.1 illustrate the robustness of the proposed bound w.r.t. different initial labelings. We initialized the figure region with (1) the whole image domain (1st row); (2) a part of the figure and ground (2nd row); and (3) a part of the figure (3rd row). Our algorithm converged to the same final solution. The typical examples in Figure 6.2 use various initializations that, by definition, preclude the application of the bound in [4]. It demonstrates how the proposed bound can deal with these different types of initializations, unlike the bound in [4], which requires the initial region to include the target region.

**6.3. Segmentation examples:** Figure 6.3 depicts some representative examples of segmentations with the proposed method along with the ground truth segmentations from the GrabCut database. The pixel error, the initial Bhattacharyya measure, the Bhattacharyya measure corresponding to the final segmentation, total computational time and image size are given for each example. The results with difficult examples undergoing strong camouflage demonstrate the effectiveness of the proposed method.
Figure 6.1. Evolution of the solution of the proposed algorithm w.r.t. different initial labelings. We initialized the figure region with (1) the whole image domain (1st row); (2) a part of the figure and ground (2nd row); and (3) a part of the figure (3rd row). The algorithm converged to the same final solution regardless of initialization. The density is estimated from the RGB components with a 3-dimensional histogram consisting of $256 \times 256 \times 256$ bins.

6.4. The effect of kernel width $\sigma$. Figure 6.4 depicts the effect of kernel width $\sigma$ on the results. This typical example shows that the histogram representation, which can be viewed as a kernel density estimate with a very small $\sigma$, yields the best performance. The higher $\sigma$, the further the result from the ground truth. This is expected in camouflage examples, where the figure and ground regions have some similarities in their color distributions. In such cases, a high value of $\sigma$ affects the ability of the algorithm as to distinguishing fine color differences.

It is worth noting that bandwidth selection for density estimation has been intensively studied in statistics [30, 23], and there are several ways of computing the optimal value of parameter $\sigma$. Generally, the optimal kernel width minimizes a given error between the original density and its estimate, and depends on the size and statistics of the data [30]. Note that more accurate estimate of the optimal kernel width can be computed via iterative procedures.
Figure 6.2. Robustness of the proposed algorithm w.r.t. initializations and comparison to BMGC [4] We initialized the figure region with (1) the whole image domain; (2) a part of the figure and ground; and (3) a part of the figure (1st row). $C(x) = 1.5$ and $c = 2$. The results demonstrate that the proposed bound (2nd row) is more robust than the bound in [4] (3rd row), which requires the initial region to include the target region. The density is estimated from the RGB components with a 3-dimensional histogram consisting of $256 \times 256 \times 256$ bins.

[30]. However, embedding such procedures into segmentation leads to a heavy computation load.
Figure 6.3. Segmentation results of the proposed method from sample images from the Grab-Cut database. 1st row: the image and ground truth segmentation boundary (green curve). 2nd row: foreground segmentation obtained with the proposed method. The density is estimated from the RGB components with a 3-dimensional histogram consisting of $256 \times 256 \times 256$ bins.

![Segmentation results](image)

<table>
<thead>
<tr>
<th>Error</th>
<th>0.02%</th>
<th>0.03%</th>
<th>0.06%</th>
<th>0.15%</th>
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<tr>
<td>$B_{\text{initial}}$</td>
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<td>$450 \times 640$</td>
<td>$600 \times 450$</td>
<td>$481 \times 321$</td>
</tr>
</tbody>
</table>

Figure 6.4. The effect of kernel width $\sigma$.

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